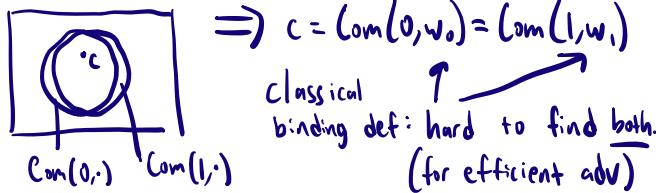


How should we define post-quantum binding?

$$\begin{array}{c} S(m) \xrightarrow{c} R \\ w \in \{0,1\}^n \\ c = \text{Com}(m, w) \end{array}$$

Open: $m, w \xrightarrow{c} R$ Verify $\text{Com}(m, w) = c$

For today: suppose Com statistically hides m



$$\begin{array}{c} \tilde{S} \xrightarrow{c} R \quad (\text{Imagine } \tilde{S} \text{ is betting on the outcome of the election}) \\ \text{After election:} \\ \text{Decide to open} \\ \text{to } m \\ (\text{the winner}) \end{array}$$

$M, W_m \xrightarrow{c} R$ Obvious; if \tilde{S} can do this, it can output w_0, w_1

\Downarrow
break underlying assumption
(last time: find collisions in CRHF)

What if \tilde{S} has a QC?

$$\tilde{S} \xrightarrow{c} R \quad (\text{Communication still classical})$$

Decide after election which m to open

Does \tilde{S} know w_0, w_1 ?

Detail: QC Recap

- QC can: apply efficient unitary U on $|0^n\rangle$. (recall: $U^\dagger U = I$)
 - measure qubits.
 - Ex: $|0^n\rangle \xrightarrow{H^{\otimes n}} \sum_{x \in \{0,1\}^n} |x\rangle |0\rangle$ any classical function f
- $$\begin{aligned} &\sum_{x \in \{0,1\}^n} |x\rangle |f(x)\rangle \\ &\downarrow \\ &\sum_{x \in \{0,1\}^n} |x\rangle |t(x)\rangle \end{aligned}$$

Hypothetical attack:

$$\begin{array}{c} \sum_{m,w} |m\rangle |w\rangle |\text{Com}(m,w)\rangle \\ \downarrow \\ |\Psi_c\rangle := \sum_{m,w} |m\rangle |w\rangle \boxed{x} = c \\ \text{both 0 and 1 in superposition} \\ \text{Com}(m,w) = c \\ \downarrow \\ |\Psi_c\rangle \xrightarrow{U_0} \boxed{x} = (0, w_0) \\ \text{OR} \\ \downarrow \\ |\Psi_c\rangle \xrightarrow{U_1} \boxed{x} = (1, w_1) \end{array}$$

Later, open to any bit m

- ① \tilde{S} successfully cheats
 - ② measurement destroys $|\Psi_c\rangle$
- $\Rightarrow \tilde{S}$ can't output both w_0, w_1 .

Summary: Binding should imply \tilde{S} can't change mind

- Classically: this means can't output w_0, w_1 .

- Quantumly: new def called collapse-binding.

(Classically, def + equivalence are immediate.
Quantumly, neither def nor equivalence are obvious)

Collapse-binding [Unruh 16]

$$\tilde{S} \xrightarrow{c} \text{Challenger}$$

- 1) Verify opening:
 $\sum_{m,w} |m\rangle |w\rangle |\text{Com}(m,w)\rangle$
 - 2) May be entangled with \tilde{S} 's state
 $\sum_{m,w} |m\rangle |w\rangle |\text{Com}(m,w)\rangle$
(if not c , abort) $\boxed{x} = c$
 - 3) $b \leftarrow \{0,1\}$
 - If $b=0$: do nothing
 - If $b=1$: measure m .
- \downarrow
- $b' \xleftarrow{\rho_{M,W}}$
- Goal: understand why this says \tilde{S} can't change its mind

For all efficient \tilde{S} :

$$\Pr[b' = b] \leq \frac{1}{2} + \text{negl}$$

Goal: understand why this says \tilde{S} can't change its mind

Quantum Information

- $\tilde{\Pi}$ is a projector if $\tilde{\Pi}^2 = \tilde{\Pi}$ and $\tilde{\Pi} = \tilde{\Pi}^\dagger$.

$$P = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}$$

- A projective measurement $\{\tilde{\Pi}_i\}_i$ satisfies $\sum_i \tilde{\Pi}_i = I$.

$$\tilde{\Pi}_1 = \begin{bmatrix} 1 & 0 \\ 0 & 0 \end{bmatrix}, \quad \tilde{\Pi}_2 = \begin{bmatrix} 0 & 0 \\ 0 & 1 \end{bmatrix}$$

If I measure $|1\rangle$, get
 $\tilde{\Pi}_1|1\rangle$ w/ prob $\|\tilde{\Pi}_1|1\rangle\|^2$.

- State after measurement is a distribution over pure states. If a system is in state $|\psi_i\rangle$ w/ prob p_i , it can be described by a density matrix

$\rho = \sum_i p_i |\psi_i\rangle \langle \psi_i|$, which captures all observable properties of the system.

- Key point: many different ensembles of states can give the same ρ .

$$\text{Ex: } \frac{1}{2}|0\rangle\langle 0| + \frac{1}{2}|1\rangle\langle 1| = \frac{1}{2}|+\rangle\langle +| + \frac{1}{2}|-\rangle\langle -| = \begin{pmatrix} 1/2 & 0 \\ 0 & 1/2 \end{pmatrix}$$

ρ is a density matrix $\Leftrightarrow \text{Tr}(\rho) = 1$ and positive semidefinite. ($\langle \psi | \rho | \psi \rangle \geq 0$ for all $|\psi\rangle$)

- Paulis: bit flip $X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$, phase flip $Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$ ($XZ = -ZX$)

1,-1 eigenvectors of X are $|+\rangle = |0\rangle + |1\rangle$, $|-\rangle = |0\rangle - |1\rangle$.
 1,-1 eigenvectors of Z are $|0\rangle, |1\rangle$

- Measure in standard basis \Leftrightarrow
 - $r \in \{0, 1\}$.
 - Apply Z^r .

density matrix: $\begin{pmatrix} a & b \\ c & d \end{pmatrix} \xrightarrow{\text{measure}} \begin{pmatrix} a & 0 \\ 0 & d \end{pmatrix} + \begin{pmatrix} 0 & b \\ c & 0 \end{pmatrix} = \begin{pmatrix} a & b \\ c & d \end{pmatrix}$
 $\downarrow \text{random } Z^r$

$$\frac{1}{2} \begin{pmatrix} a & b \\ c & d \end{pmatrix} + \frac{1}{2} \begin{pmatrix} a & -b \\ -c & d \end{pmatrix} = \begin{pmatrix} a & 0 \\ c & d \end{pmatrix}$$

- Overlap between two states: $|0\rangle$ and $|1\rangle$: Prepare $\frac{1}{\sqrt{2}}(|0\rangle|0\rangle + |1\rangle|1\rangle)$, measure first qubit in $|+\rangle, |-\rangle$ basis.

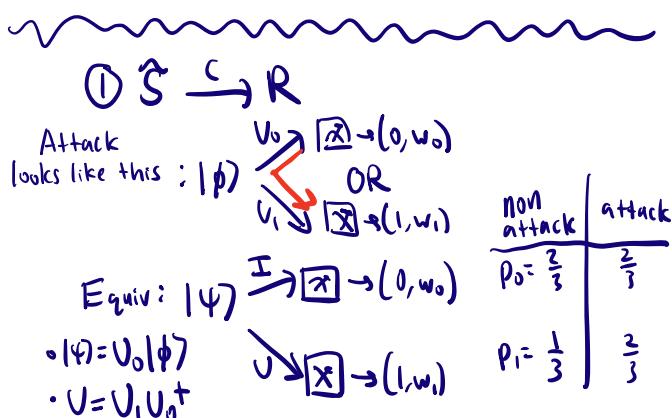
Intuition: if $|0\rangle$ and $|1\rangle$ are parallel, get interference, if perpendicular, no interference.

$$\text{If } |0\rangle|0\rangle + |1\rangle|1\rangle = |+\rangle|+\rangle \Rightarrow \text{always get } |+\rangle$$

$$\text{If } |0\rangle|0\rangle + |1\rangle|1\rangle^\dagger \Rightarrow \frac{1}{2} \text{ prob of } |+\rangle.$$

In general: Pr of + on $|0\rangle|0\rangle + |1\rangle|1\rangle$ is $\frac{1}{2} + \text{Re}(\langle \tilde{\Phi}_0 | \tilde{\Phi}_1 \rangle)$.

Goal: ① What does "changing your mind" mean?
 ② collapse-binding \Rightarrow can't change mind



Define projection $\tilde{\Pi}_0$ onto valid 0 openings w_0 .

$$\tilde{\Pi}_0 = |0\rangle\langle 0|_M \otimes \sum_{w_0} |w_0\rangle\langle w_0| \quad (\text{define } \tilde{\Pi}_1 \text{ similarly})$$

Must guarantee:

- Adv can't open in two ways

$$\|\tilde{\Pi}_0|\psi\rangle\|^2 + \|\tilde{\Pi}_1|\psi\rangle\|^2 = p_0 + p_1 = \text{negl}$$

- Adv can't change 0 opening w_0 into 1 opening w_1 :

$$\|\tilde{\Pi}_1 V \tilde{\Pi}_0|\psi\rangle\|^2 = \text{negl}.$$

Lemma:

$$\|\tilde{\Pi}_1 V \tilde{\Pi}_0|\psi\rangle\|^2 = \varepsilon \Rightarrow \|\tilde{\Pi}_0|\psi\rangle\|^2 + \|\tilde{\Pi}_1 V |\psi\rangle\|^2 \leq 1 + 2\varepsilon + \varepsilon^2$$

(We'll prove this later).

Can't change mind \Rightarrow Can't open in two ways.

suffices to show C.B. $\Rightarrow \langle \psi | \tilde{\Pi}_0 V^\dagger \tilde{\Pi}_1 V \tilde{\Pi}_0 |\psi\rangle = \text{negl}$

②

$$\langle \tilde{\Phi}_0 | \tilde{\Phi}_1 \rangle$$

Why does $\langle \tilde{\Phi}_0 | \tilde{\Pi}_0 V^\dagger \tilde{\Pi}_1 V \tilde{\Pi}_0 | \tilde{\Phi}_1 \rangle = \varepsilon$ imply distinguisher for C.B.?

Intuition: Want this inner product to show up distinguishing probability.

Observation: This inner product corresponds to our ability to distinguish $|0\rangle|\tilde{\Phi}_0\rangle + |1\rangle V^\dagger |\tilde{\Phi}_1\rangle = |\tilde{\Phi}_0\rangle$ from $|0\rangle|\tilde{\Phi}_0\rangle - |1\rangle V^\dagger |\tilde{\Phi}_1\rangle = |\tilde{\Phi}_1\rangle$

measure first qubit in +,- basis, guess 0 if +, 1 if -.

$$\Pr[\text{guess } b] - \frac{1}{2} = \frac{1}{2} [\text{measure + on } |\tilde{\Phi}_0\rangle] - \frac{1}{2} [\text{measure + on } |\tilde{\Phi}_1\rangle]$$

$$= \frac{1}{2} \left(\frac{1}{2} + \langle \tilde{\Phi}_0 | V^\dagger | \tilde{\Phi}_1 \rangle \right) - \frac{1}{2} \left(\frac{1}{2} - \langle \tilde{\Phi}_0 | V^\dagger | \tilde{\Phi}_1 \rangle \right)$$

$$= \langle \tilde{\Phi}_0 | V^\dagger | \tilde{\Phi}_1 \rangle = \varepsilon.$$

So if $\langle \tilde{\phi}_0 | U^+ | \tilde{\phi}_1 \rangle = \epsilon$, I can detect a phase flip. Intuitively, C.B. is also about detecting a phase flip, since challenger either applies I or Z .

- 1) Construct $|0\rangle|\tilde{\phi}_0\rangle + |1\rangle|\tilde{\phi}_1\rangle$ (see below)
- 2) Send M,W registers of state to challenger
 - If $b=0$: get $|0\rangle|\tilde{\phi}_0\rangle + |1\rangle|\tilde{\phi}_1\rangle$.
 - If $b=1$: get $|0\rangle(Z_M \otimes I)|\tilde{\phi}_0\rangle + |1\rangle(Z_M \otimes I)|\tilde{\phi}_1\rangle$
 $= |0\rangle|\tilde{\phi}_0\rangle - |1\rangle|\tilde{\phi}_1\rangle$.
- 3) do $CNOT-V^+$, measure first qubit in $+/-$ basis.
 Guess $b' = 0$ if $+$, $b' = 1$ if $-$.
 Then $\Pr[b' = b] = \frac{1}{2} + \epsilon$.

Aside: How we prepare $|0\rangle|\tilde{\phi}_0\rangle + |1\rangle|\tilde{\phi}_1\rangle$

- Measure $|4\rangle$ w/ $\{T_{10}, I - T_{10}\}$.
 (If $I - T_{10}$ outcome, give up)
- Prepare ancilla $|+\rangle$.
- Do controlled-V on $|+\rangle T_{10} |4\rangle$
 to get $|0\rangle T_{10} |4\rangle + |1\rangle V T_{10} |4\rangle$.
- Measure w/ $T_1' = (X \otimes I + I \otimes X) \otimes T_{10}$, get
 $|0\rangle T_{10} |4\rangle + |1\rangle T_1 V T_{10} |4\rangle$

Want: Completeness: Honest P always convinces V.

Soundness: If G has no H-cycle, then V accepts w/ prob $\leq \frac{1}{2} + \text{negl}$, even if \tilde{P} is malicious

Zero-Knowledge: \tilde{V} learns nothing
(not today) (beyond the fact that G has an H-cycle)

Known: classical-binding commitments



Blum is sound against classical \tilde{P}

W.T.S.: collapse-binding commitments



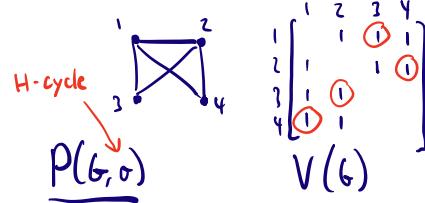
Blum is sound against quantum \tilde{P}

So far, we've seen why C.B. captures PQ binding. Today: How do we use C.B. commitments to build PQ secure protocols?

Next time: constructing C.B. commitments

Blum's Hamiltonian Cycle Protocol

Goal: P wants to convince V that G has an H-cycle or without revealing σ .



- $\pi \leftarrow S_n$ (permute vertices)
 Commit to each bit of $\pi(G), \pi_i$.
 $\pi(G), \pi_i \xrightarrow{\{C_i\}_i}$
 $b \leftarrow \{0, 1\}$
 - If $b=0$: open everything
 $z_0 = (\pi(G), \pi_i, \text{all } w_i)$
 - If $b=1$: only open
 $\pi_i(\sigma)$ commitments
 $z_1 = (\pi(\sigma), \text{corresponding } w_i)$
- Accept if:
 $\begin{cases} b=0: \text{valid perm of } G \text{ and openings verify} \\ b=1: \text{valid cycle and openings verify} \end{cases}$

Easy observation:

- c is "good" if

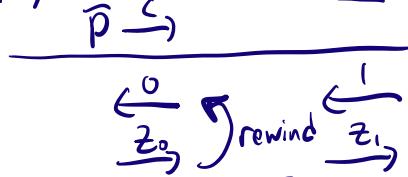
\tilde{P} convinces V
w/ prob $\geq \frac{1}{2} + \frac{\epsilon}{2}$
conditioned on c.

\tilde{P} sends good c w/ prob $\geq \frac{\epsilon}{2}$.

(Pf: Suppose not. Then \tilde{P} wins w/ prob at most
 $(1 - \frac{\epsilon}{2})(\frac{1}{2} + \frac{\epsilon}{2}) + \frac{\epsilon}{2} = \frac{1}{2} + \epsilon - \frac{\epsilon^2}{4} < \frac{1}{2} + \epsilon$.)
Contradiction

Rest of today: assume \tilde{P} wins w/ $\frac{1}{2} + \frac{\epsilon}{2}$
prob on fixed c

Key Idea: Run \tilde{P} on both $b=0, 1$.



Let $p_b = \Pr_{P(b) \rightarrow z_b}[z_b \text{ valid}]$. Know: $p_0 + p_1 = 1 + \epsilon$

$$\Pr[z_0, z_1 \text{ both valid}] \geq 1 - \Pr[z_0 \text{ invalid}] - \Pr[z_1 \text{ invalid}] \\ = 1 - (1 - p_0) - (1 - p_1) = p_0 + p_1 - 1 = \epsilon$$

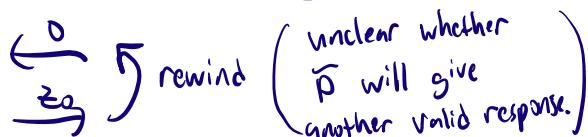
Last step: if z_0, z_1 both valid, we get H-cycle.

- z_0 gives $\pi(G), \pi$. z_1 gives τ .
- If τ consistent w/ $\pi(G)$, then $\pi^{-1}(\tau)$ is H-cycle of G
- If τ inconsistent w/ $\pi(G)$, then some commitment must be opened to both 0, 1.
 \Rightarrow break classical binding.

PQ Soundness of Blum

Difficulty: if \tilde{P} is quantum, recording z_0 can disturb its state.

$\tilde{P} \xrightarrow{C}$



We'll show: if commitments are C.B., possible to make rewinding work!

Sub-claim: Ext' succeeds w/ prob $\Omega(\varepsilon^2)$.

Same Lemma from earlier:

$$\text{If } \underbrace{\|\Pi_0^V|4\rangle\|^2}_{p_0} + \underbrace{\|\Pi_1^V|4\rangle\|^2}_{p_1} \geq 1 + \varepsilon, \\ \text{then } \|\Pi_1^V \sqrt{p_0} \Pi_0^V|4\rangle\|^2 \geq \Omega(\varepsilon^2).$$

Pf of Claim 2

If Ext succeeds, why does it output an H-cycle?

- Recall: If τ and $\pi(G)$ are consistent, then $\pi^{-1}(\tau)$ is H-cycle of G .
- C.B. implies adv can't change mind!

$\|\Pi_1^V \sqrt{p_0} \Pi_0^V|4\rangle\| = \text{negl}$, so we'll never see a 0 in $\pi(G)$ turn into a 1 in τ

(other direction also true, but remember that V rejects if τ has any 0's)

\tilde{P} has state $|4\rangle_{R,H}$ after sending C .

WLOG (can always redefine $\pi(G), V$ to ensure this)

- To output z_0 , it measures R register.
- To output z_1 , it applies V , then measures R .

Let Π_b^V be projection onto accepting z_b ,

$$\Pi_b^V := \sum_{z_b: V_{c,b}(z_b) = \text{acc}} |z_b\rangle\langle z_b|.$$

$$\text{Assumption: } p_0 + p_1 \geq 1 + \varepsilon, \text{ i.e., } \underbrace{\|\Pi_0^V|4\rangle\|^2}_{p_0} + \underbrace{\|\Pi_1^V|4\rangle\|^2}_{p_1} \geq 1 + \varepsilon$$

$\text{Ext}(V, |4\rangle_{RH})$:

- Measure $V_{c,0}(z_0) \rightarrow \text{acc/rej}$ (if rej, stop/fail)
- Measure $\pi(G), \pi$ (part of z_0)
- Apply V .
- Measure $V_{c,1}(z_1) \rightarrow \text{acc/rej}$ (if rej, stop/fail)
- If successful, measure τ , output $\pi^{-1}(\tau)$.

W.t.s: Claim 1: Ext succeeds w/ prob poly(ε).

Claim 2: If Ext succeeds, it outputs H-cycle.

Pf of Claim 1

Consider Ext' that skips (2).

$$|\Pr[\text{Ext}' \text{ succeeds}] - \Pr[\text{Ext} \text{ succeeds}]| = \text{negl}.$$

If not, get efficient distinguisher against C.B.

Missing Lemma:

For projectors Π_0, Π_1 , state $|4\rangle$, if $\|\Pi_1 \Pi_0|4\rangle\| = \varepsilon$, then $\|\Pi_0|4\rangle\|^2 + \|\Pi_1|4\rangle\|^2 \leq 1 + 2\varepsilon + \varepsilon^2$.

(Looks different, but we can set $\Pi_0 = \Pi_0^V$ and $\Pi_1 = V^* \Pi_1^V V$)

$$\text{Pf: } \|\Pi_0|4\rangle\|^2 + \|\Pi_1|4\rangle\|^2$$

$$= \|\Pi_0|4\rangle\|^2 + \|\Pi_1 \Pi_0|4\rangle + \Pi_1(I - \Pi_0)|4\rangle\|^2$$

triangle inequality

$$\leq \|\Pi_0|4\rangle\|^2 + (\|\Pi_1 \Pi_0|4\rangle\| + \|\Pi_1(I - \Pi_0)|4\rangle\|)^2$$

$$= \|\Pi_0|4\rangle\|^2 + \varepsilon^2 + 2\varepsilon \underbrace{\|\Pi_1(I - \Pi_0)|4\rangle\|}_{\leq 2\varepsilon} + \|\Pi_1(I - \Pi_0)|4\rangle\|^2$$

$$\leq \|\Pi_0|4\rangle\|^2 + \|(I - \Pi_0)|4\rangle\|^2 + \varepsilon^2 + 2\varepsilon = 1 + \varepsilon^2 + 2\varepsilon.$$