Quantum Commitments and Black Hole Radiation Decoding

Fermi Ma (Simons and Berkeley)

Based on discussions with Sam Gunn (Berkeley) and Alex Lombardi (Berkeley → Princeton)

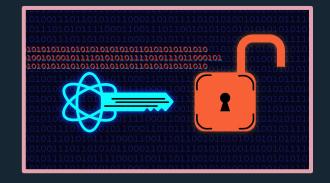
Question: What does black hole radiation decoding have to do with quantum cryptography?





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Answer:

Black-Hole Radiation Decoding is Quantum Cryptography

Zvika Brakerski*

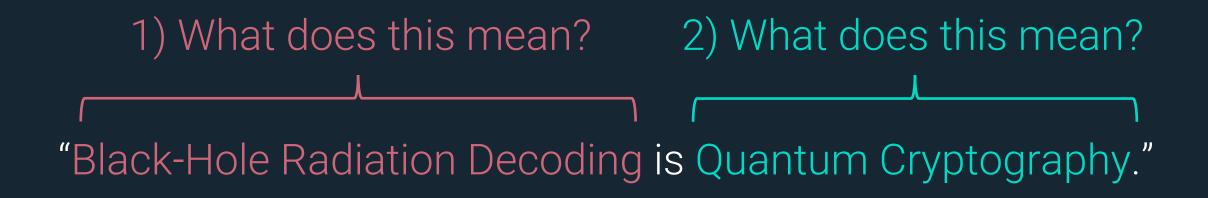
Abstract

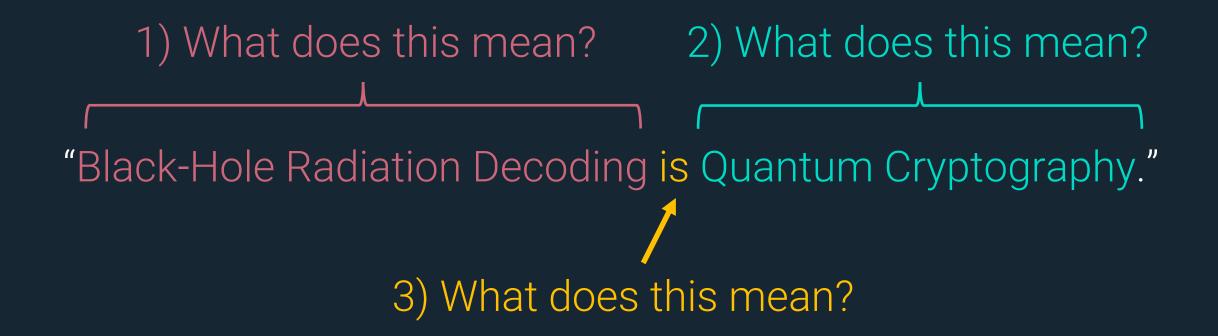
We propose to study equivalence relations between phenomena in high-energy physics and the existence of standard cryptographic primitives, and show the first example where such an

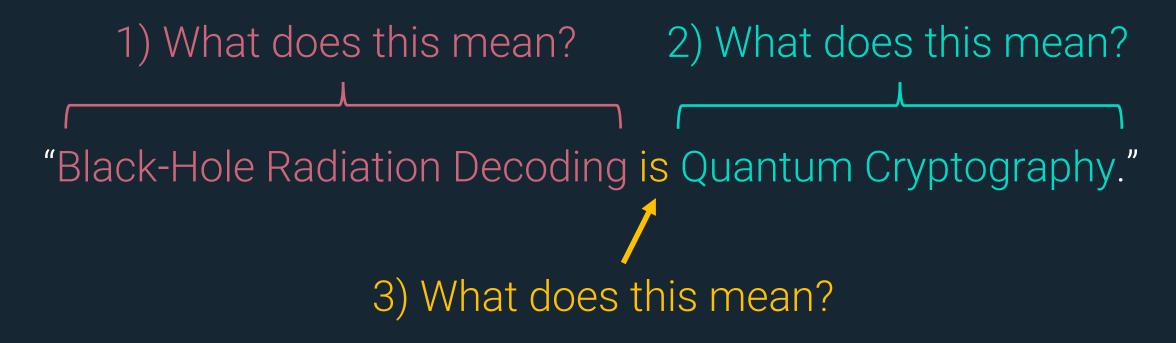
(building on [Harlow-Hayden13, Aaronson16])

"Black-Hole Radiation Decoding is Quantum Cryptography."

1) What does this mean? "Black-Hole Radiation Decoding is Quantum Cryptography."







Goal: understand the title of Zvika's paper

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Warning: I'm not a physicist.

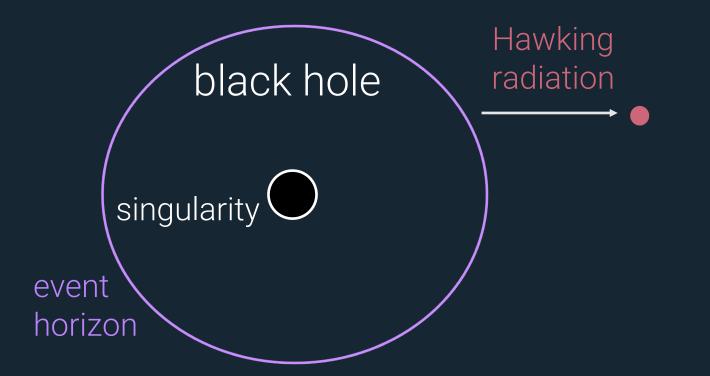


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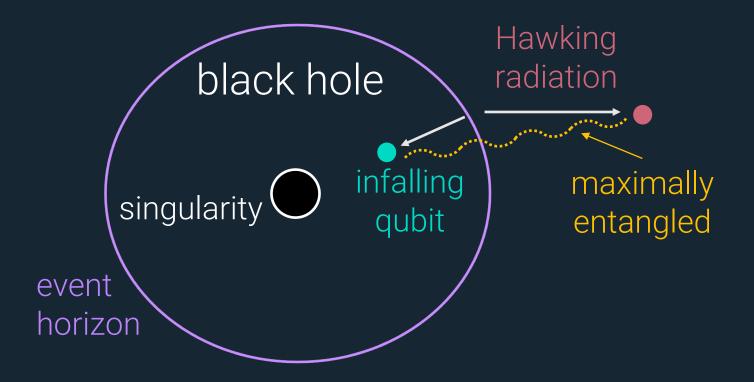
Everything I'm about to say about black hole physics is from Scott Aaronson's Barbados lecture notes (any mistakes are my own).



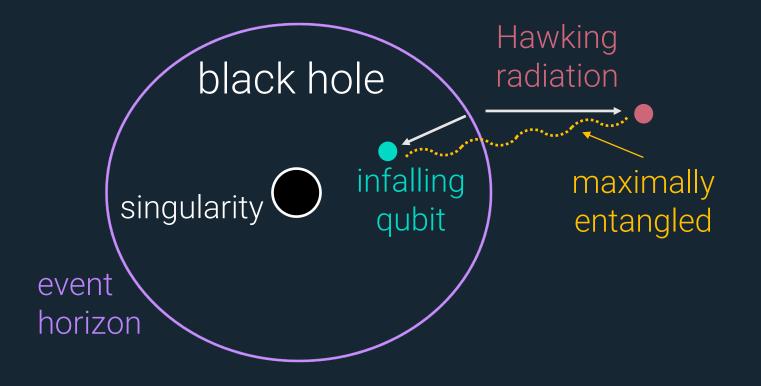
• Black holes emit qubits of Hawking radiation.

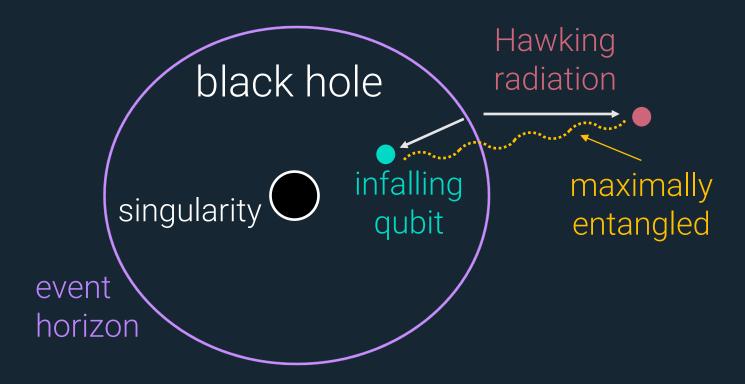


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- Each outgoing qubit is maximally entangled with an infalling qubit.

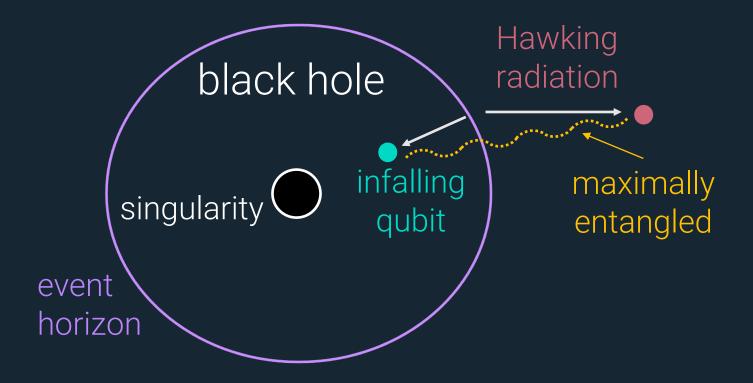


- Black holes emit qubits of Hawking radiation.
- Each outgoing qubit is maximally entangled with an infalling qubit.
- After long enough, black hole evaporates completely.

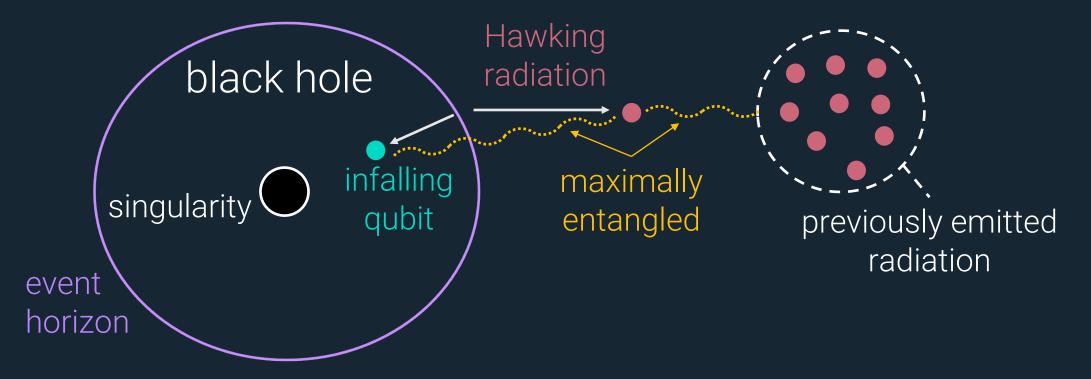




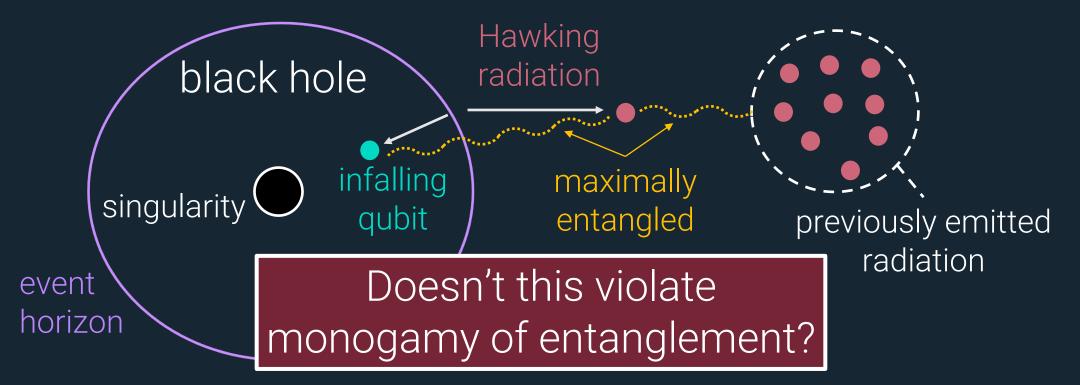
• Post-evaporation state is a (roughly) a random pure state.

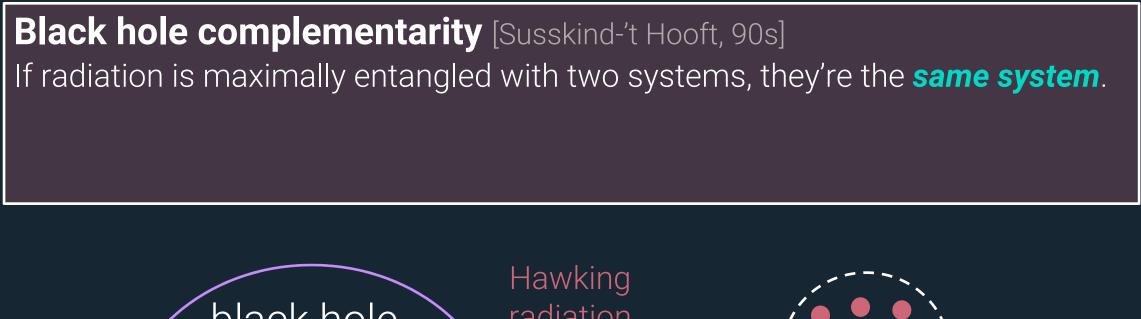


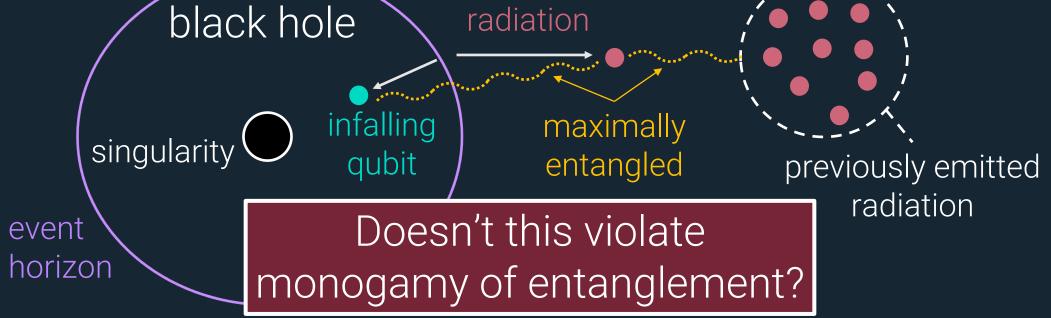
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- Consequence: after ~1/2 of the black hole has evaporated, outgoing qubits are maximally entangled with previously emitted radiation.



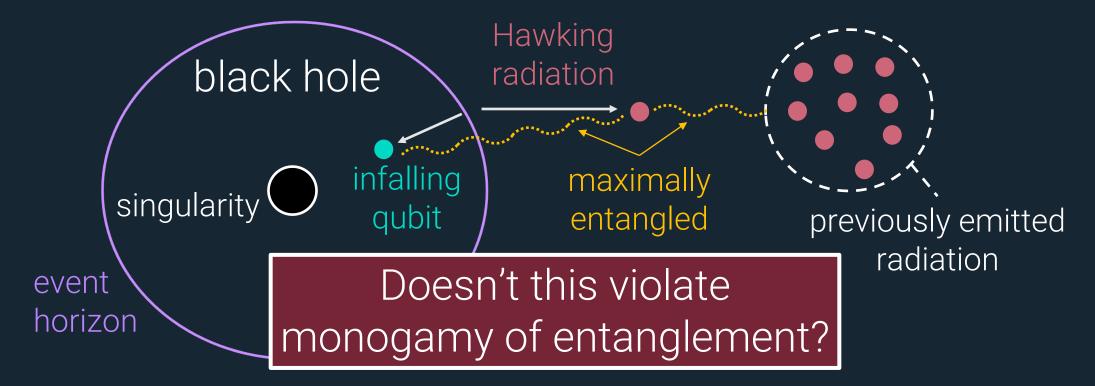
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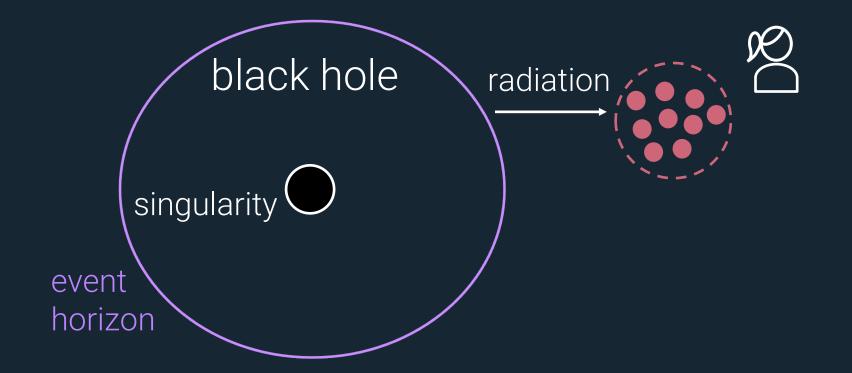
Black hole complementarity [Susskind-'t Hooft, 90s]
If radiation is maximally entangled with two systems, they're the same system.
Firewall paradox [Almheiri-Marolf-Polchinski-Sully, 11]
Thought experiment in which an observer detects the monogamy violation.



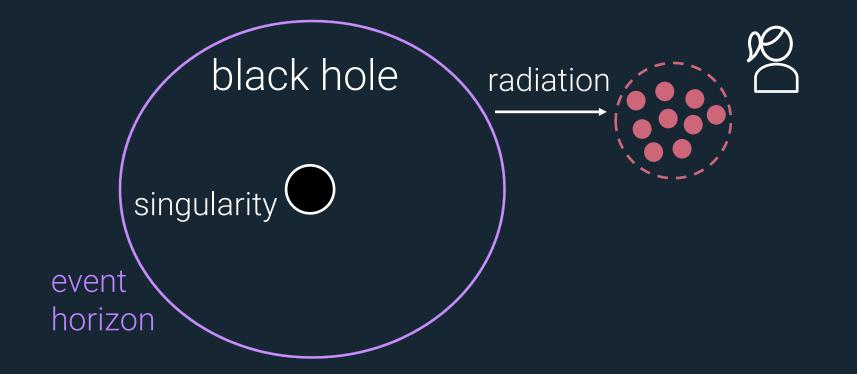




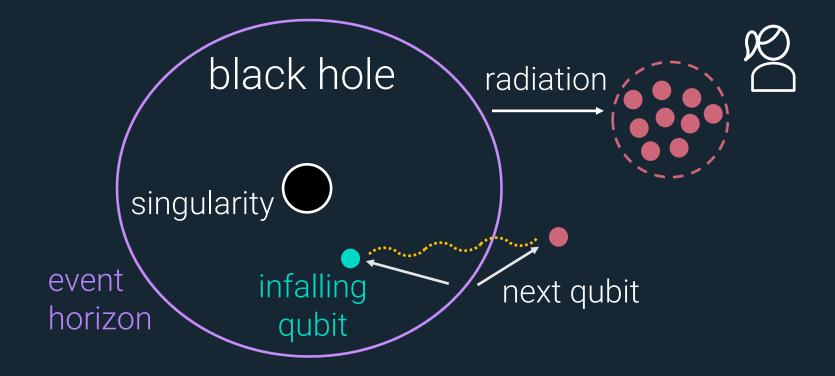
1) Alice collects radiation until 2/3 of black hole has evaporated.



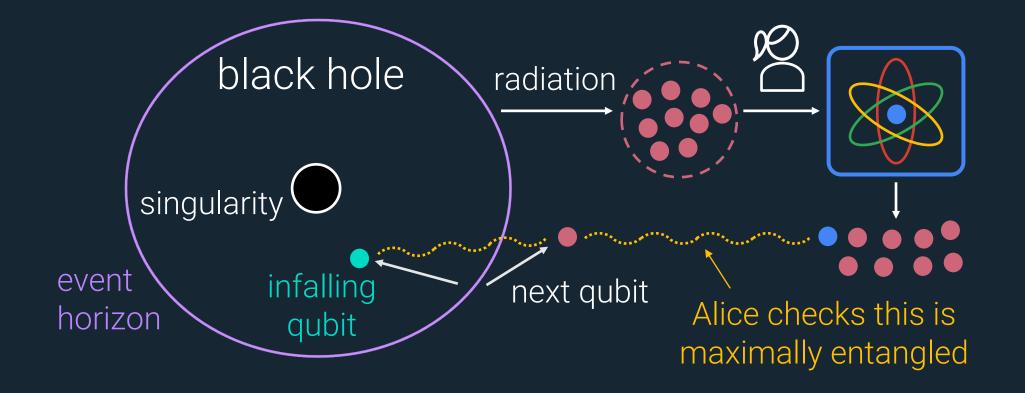
 Alice collects radiation until 2/3 of black hole has evaporated.
 Alice uses a quantum computer to "check" that the next qubit is entangled with her collected radiation (e.g., distills an EPR pair).



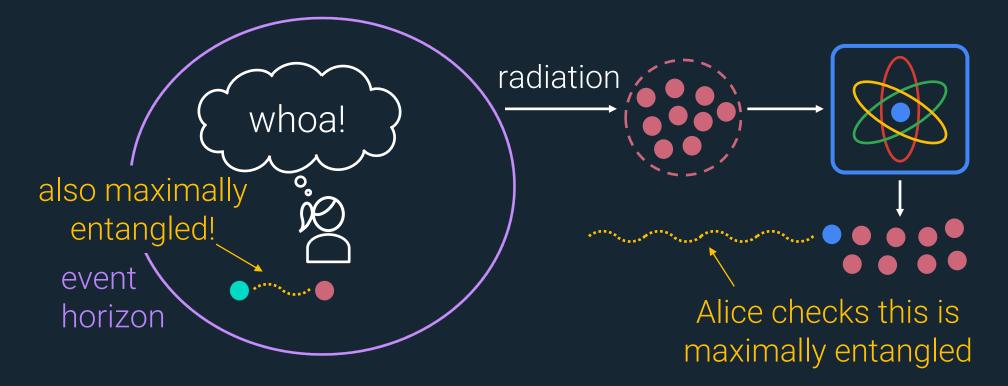
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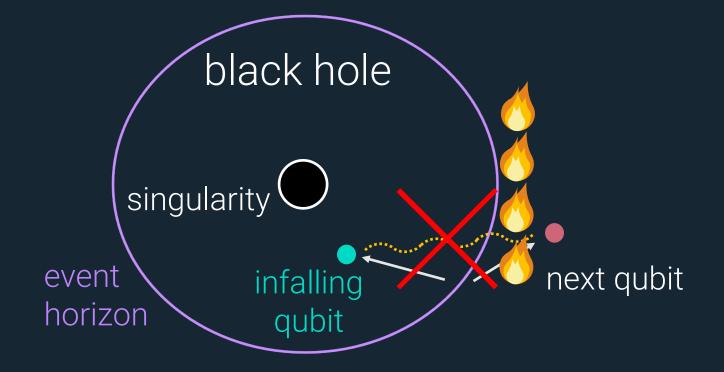


 Alice collects radiation until 2/3 of black hole has evaporated.
 Alice uses a quantum computer to "check" that the next qubit is entangled with her collected radiation (e.g., distills an EPR pair).
 Alice jumps into the black hole.



AMPS11 proposed resolution:

"Firewall" outside event horizon (breaking entanglement)



In 2013, Harlow and Hayden proposed a different resolution to the AMPS paradox based on **computational complexity**.

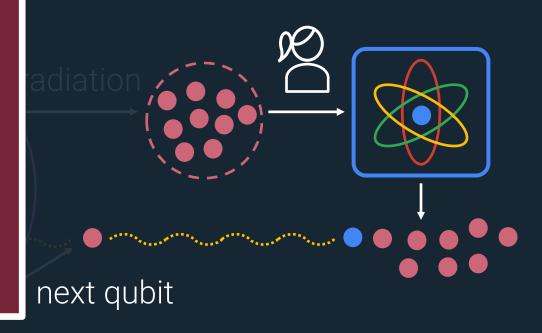
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Very cool and surprising!!

(AMPS11) experiment:
(Alice collects radiation until 2/3 of black hole has evaporated.
2) Alice uses a quantum computer to "check" that the next qubit is entangled with her collected radiation (e.g., distills an EPR pair).

[Harlow-Hayden 2013]

Under certain cryptographic assumptions, this step can require **exponential** time.

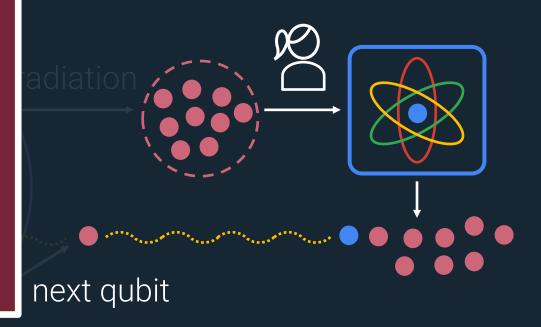


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[Harlow-Hayden 2013]

Under certain cryptographic assumptions, this step can require **exponential** time.

By the time she's done decoding, the black hole will have evaporated!



Plan for this talk

(1) Background on black holes

(2) Radiation decoding problem [Harlow-Hayden13]

(3) Radiation *distinguishing* problem

(4) Connection to quantum commitments

(3) + (4) is an alternative view of [Brakerski23].

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Task: Given **R** register of $|\psi\rangle_{\text{RBH}} = C|0^n\rangle$, output a single qubit **A** such that (**A**, **B**) is the EPR state $|00\rangle + |11\rangle$.

(promised that **R** and **B** are maximally entangled)

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"Hard" means no QPT adversary can win with probability $\geq \frac{1}{4} + \operatorname{negl}(n)$ (formalized by [Brakerski23])

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Later works *weakened* the assumptions needed:

- [Aaronson16]: quantum-secure one-way functions
- [Brakerski23]: quantum bit commitment

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"This can be viewed (with proper disclaimers, as we discuss) as providing a physical justification for the existence of secure cryptography" – [Brakerski23]

Rest of today: new perspective on Brakerski's result/proof.

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(2) Radiation decoding problem [Harlow-Hayden13]

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Instead of studying the [HH13] radiation **decoding** problem, we'll define a new radiation **distinguishing** problem.

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The point:

R and **B** are maximally entangled, but this entanglement isn't efficiently detectable.

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Claim 1: Distinguishing is easier than decoding.

If you can solve the decoding problem with advantage $1/4 + \varepsilon$, you can distinguish with advantage ε .

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Claim 2: Distinguishing should still be hard.

If Alice can't trigger a firewall, then she shouldn't be able to detect entanglement between B and R in the AMPS experiment.

Plan for this talk

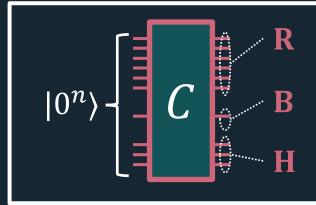
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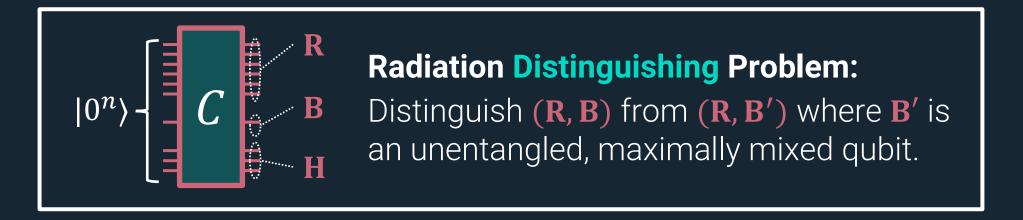
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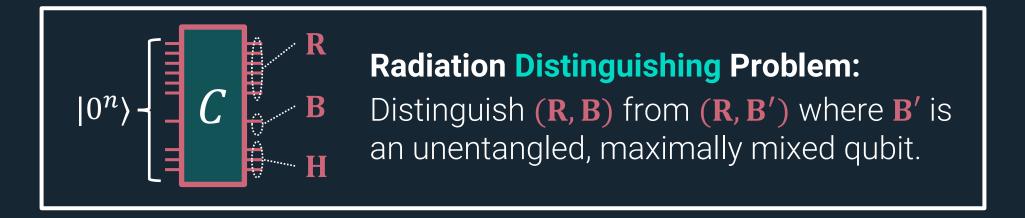


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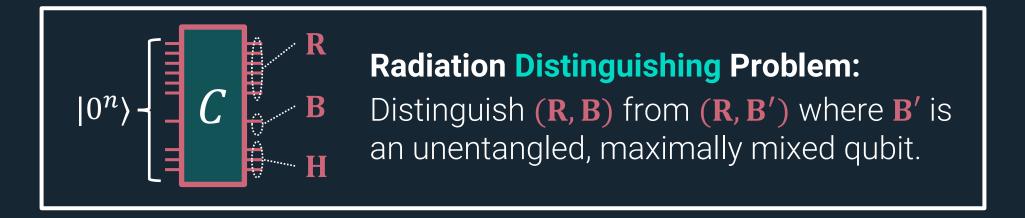


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Radiation Distinguishing is hard if and only if *quantum commitments to the EPR state* exist.

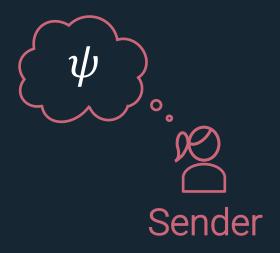


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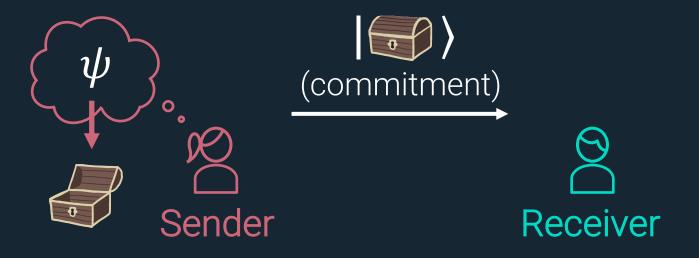
Up next: define commitments to quantum states

[Gunn-Ju-M-Zhandry23]

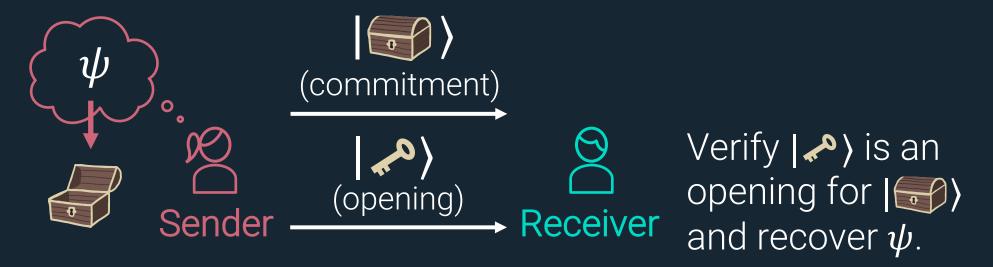




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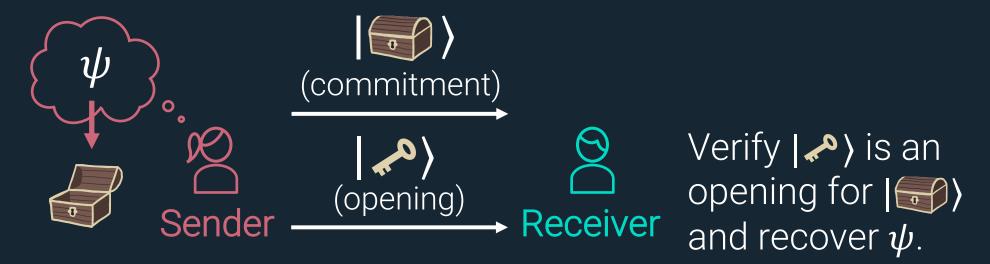


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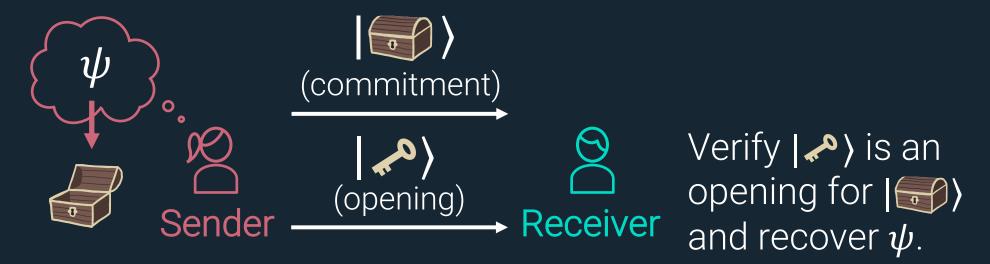
Protocol that lets a sender commit to a (possibly entangled) quantum state ψ , with the ability to reveal ψ later.



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- Requires computational assumptions [M96, LC96].
- Exist if and only if quantum bit commitments exist.

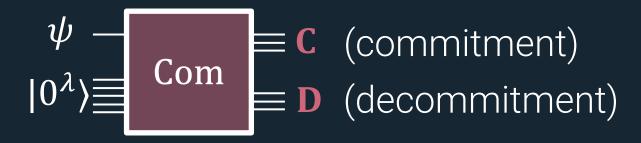
Commitment Syntax







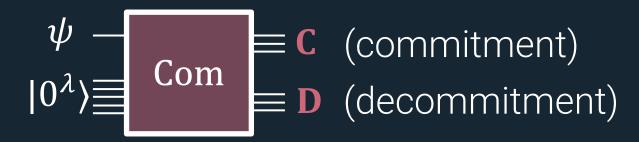
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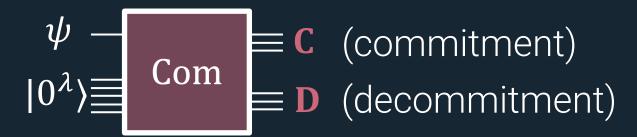


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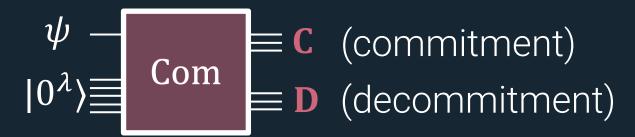


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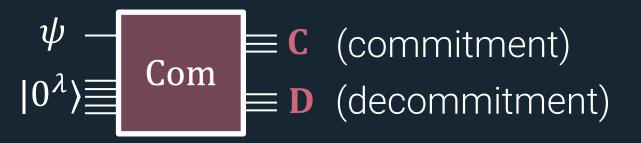


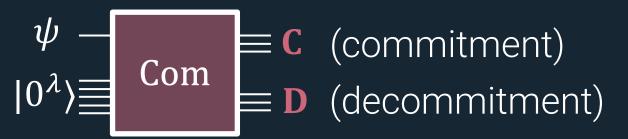
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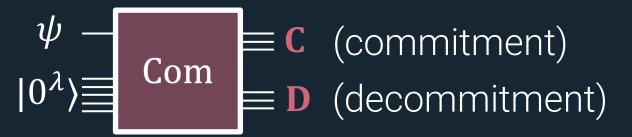


To verify (C, D), receiver applies Com^{\dagger} and checks if last λ bits are 0.





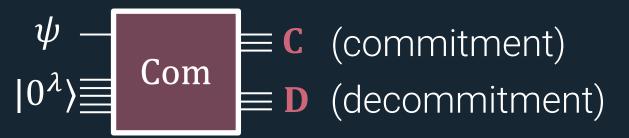
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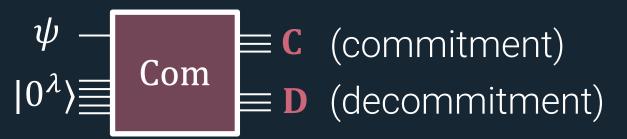
Exists an inefficient unitary $U_{\mathbf{C}}$ that recovers ψ from **C** alone.

$$\begin{array}{c} \psi \\ \psi \\ 0^{\lambda} \end{array} \end{array} = \begin{array}{c} \mathbf{C} \\ \mathbf{C} \\ \mathbf{D} \end{array} = \begin{array}{c} \mathbf{C} \\ \mathbf{D} \\ \mathbf{D} \end{array}$$



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Computational hiding: no QPT adversary can distinguish: (1) commitment to ψ of the adversary's choice (2) commitment to junk (e.g., maximally mixed state)



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Crucial point: since adversary picks ψ , indistinguishability holds even if the adversary has a state entangled with ψ .

Setup: Prepare |EPR>_{AB} and commit to **A**.

$$|EPR\rangle - \begin{cases} B & B \\ A & B \\ 0^{\lambda} \rangle \end{cases} B - B \\ Com & C \\ 0^{\lambda} \rangle B \\ 0^{\lambda} \rangle B$$

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Statistical Binding: B and C are maximally entangled.

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Computational Hiding: (**B**, **C**) indistinguishable from (**B**, **C'**) where **C'** is a commitment to a maximally mixed state

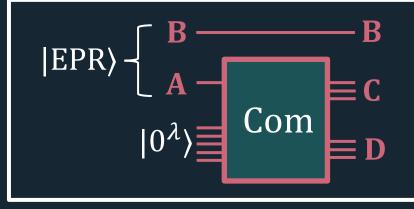
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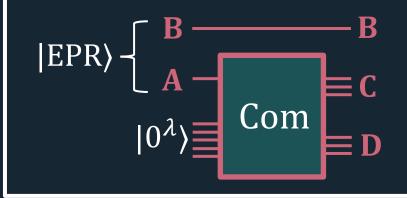
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Fact: (**B**, **C**') is distributed as (**B**', **C**) for **B**' maximally mixed.

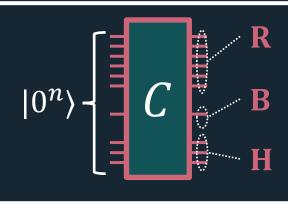


Breaking Hiding of EPR Commitment:

Promised that **B** and **C** are maximally entangled, distinguish (**B**, **C**) from (**B'**, **C**) where **B'** is an unentangled, maximally mixed qubit.



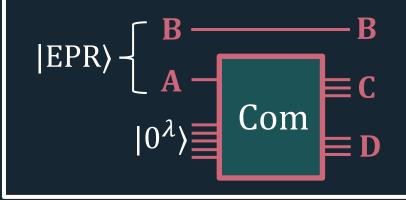
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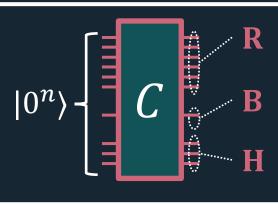
Radiation Distinguishing Problem:

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Thus, quantum commitments \rightarrow hard radiation distinguishing.



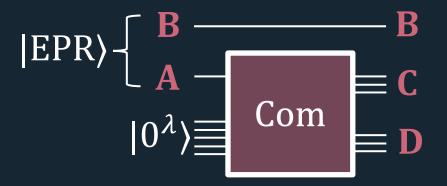
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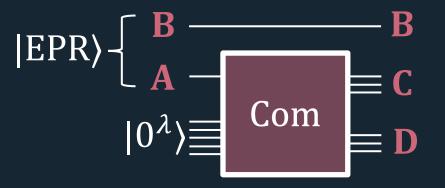


Radiation Distinguishing Problem:

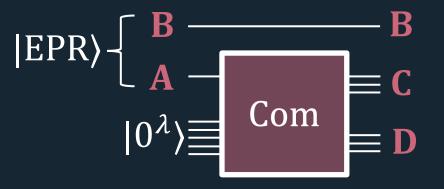
Promised that **B** and **R** are maximally entangled, distinguish (**B**, **R**) from (**B'**, **R**) where **B'** is an unentangled, maximally mixed qubit.

One last thing: to show hard radiation distinguishing \rightarrow crypto, need to show EPR commitments \rightarrow commitments to any state.



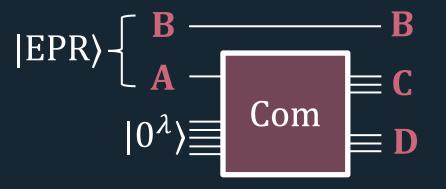


Just teleport ψ into **C**: to commit to ψ , measure (ψ , **B**) in the Bell basis to get classical bits (x, z), and send (**C**, x, z).



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- Statistical Binding: C determines A. (A, x, z) determines ψ .
- Computational Hiding: (C, x, z) indistinguishable from (C', x, z) where C' is a commitment to junk, but this is independent of ψ .

Conclusion

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Future research direction: give more evidence for hardness.

Given description of a random circuit C, how hard is it to distinguish $C|0^n\rangle$ from $C|1^n\rangle$ given 2n/3 of the qubits?